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MENTATION PAGE

Form Approved OPM No. 0704-0188



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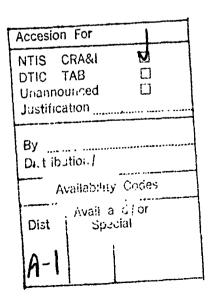
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4 TITLE AND SUBTITLE			5 FUNDING NUMBERS	
	XX/VMS=> Intel 386/WITEK 3	167 VMS 5.2	5 FORDING NUMBERS	
•), MicroVAX 3100 (Host) to iSi			
WEITEK 3167 fpu)(Target), 9		30 300/110 (using a		
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6. AUTHOR(S)	-			
Wright-Patterson AFB, Dayton	, OH]
USA				
7 PERFORMING ORGANIZATION NAME			8. PERFORMING ORGAN	ZATION
Ada Validation Facility, Langua	ige Control Facility ASD/SCEL		REPORT NUMBER	
Bldg. 676, Rm 135			AVF-VSR-404-0891	
Wright-Patterson AFB			1	
Dayton, OH 45433				
9 SPONSORING/MONITORING AGEN	ICY NAME(S) AND ADDRESS(ES)	-	10. SPONSORING/MONIT	ORING AGENCY
Ada Joint Program Office			REPORT NUMBER	
United States Department of D	lefense			
Pentagon, Rm 3E114	3.3.13.0			
Washington, D.C. 20301-3081				
				
11. SUPPLEMENTARY NOTES				
12a. DISTRIBUTION/AVAILABILITY ST	ATEMENT		12b DISTRIBUTION COD	E
Approved for public release; distribution unlimited.				
13 ABSTRACT (Maximum 200 words)				
Verdix Corporation, VADS VAX	X/VMS=> Intel 386/WFITEK 3	167. VMS 5.2. VAda	-110-03315. Version 6.0.	
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AVF Control Number: AVF-VSR-404-0891 1 August 1991 90-09-25-VRX

Ada COMPILER
VALIDATION SUMMARY REPORT:
Certificate Number: 901129W1.11094
Verdix Corporation

VADS VAX/VMS => Intel 386/WEITEK 3167, VMS 5.2, VAda-110-03315, Version 6.0 MicroVAX 3100 => iSBC 386/116 (using a WEITEK 3167 fpu)

Prepared By:
Ada Validation Facility
ASD/SCEL
Wright-Patterson AFB OH 45433-6503





Certificate Information

The following Ada implementation was tested and determined to pass ACVC 1.11. Testing was completed on 29 November 1990.

Compiler Name and Version: VADS VAX/VMS => Intel 386/WEITEK 3167,

VMS 5.2, VAda-110-03315, Version 6.0

Host Computer System:

MicroVAX 3100, VMS 5.2

Target Computer System:

Intel iSBC 386/116 (using a WEITEK 3167 fpu),

(bare machine)

Customer Agreement Number: 90-09-25-VRX

See Section 3.1 for any additional information about the testing environment.

As a result of this validation effort, Validation Certificate 901129W1.11094 is awarded to Verdix Corporation. This certificate expires on 1 March 1993.

This report has been reviewed and is approved.

Ada Validation Facility

Steven P. Wilson Technical Director

. ASD/SCEL

Wright-Patterson AFB OH 45433-6503

Ada Validation Organization
Director, Computer & Software Engineering Division
Institute for Defense Analyses

Alexandria VA 22311

Ada Joint Program Office

Dr. John Solomond, Director

Department of Defense Washington DC 20301

DECLARATION OF CONFORMANCE

Customer:

Verdix Corporation

Ada Validation Facility:

ASD/SCEL Wright Patterson AFB, OH 45433-6503

Ada Compiler Validation Capability (ACVC)

Version: 1.11

Ada Implementation

Compiler Name: VADS VAX/VMS=>Intel 386, WEITEK 3167, VMS 5.2, VAda-110-03315, Version: 6.0

Host Computer System: Micro VAX 3100, VMS 5.2

Target Computer System: iSBC 386/116 (using a WEITEK 3167 fpu), (bare machine)

Declaration

I, the undersigned, representing Verdix Corporation, declare that Verdix Corporation has no knowledge of deliberate deviations from the Ada Language Standard ANSI/MIL-STD-1815A in the implementation listed in this declaration. I declare that the Verdix Corporation is the owner of the above implementation and the certificates shall be awarded in the name of Verdix Corporation.

Date: 2/1/9/

Stechen Zangier

1600 N W Compton Drive

Suite 357

Beaverton Oregon 97006

CHAPTER 2

IMPLEMENTATION DEPENDENCIES

2.1 WITHDRAWN TESTS

The following tests have been withdrawn by the AVO. The rationale for withdrawing each test is available from either the AVO or the AVF. The publication date for this list of withdrawn tests is 12 October 1990.

E28005C	B28006C	C34006D	B41308B	C43004A	C45114A
C45346A	C45612B	C45651A	C46022A	B49008A	A74006A
C74308A	B83022B	B83022H	B83025B	B83025D	B83026B
B85001L	C83026A	C83041A	C97116A	C98003B	BA2011A
CB7001A	CB7001B	CB7004A	CC1223A	BC1226A	CC1226B
BC3009B	BD1B02B	BD1B06A	AD1B08A	BD2AO2A	CD2A21E
CD2A23E	CD2A32A	CD2A41A	CD2A41E	CD2A87A	CD2B15C
BD3006A	BD4008A	CD4022A	CD4022D	CD4024B	CD4024C
CD4024D	CD4031A	CD4051D	CD5111A	CD7004C	ED7005D
CD7005E	AD7006A	CD7006E	AD7201A	AD7201E	CD7204B
BD8002A	BD8004C	CD9005A	CD9005B	CDA201E	CE2107I
CE2117A	CE2117B	CE2119B	CE2205B	CE2405A	CE3111C
CE3118A	CE3411B	CE3412B	CE3607B	CE3607C	CE3607D
CE3812A	CE3814A	CE3902B			

2.2 INAPPLICABLE TESTS

A test is inapplicable if it contains test objectives which are irrelevant for a given Ada implementation. Reasons for a test's inapplicability may be supported by documents issued by the ISO and the AJPO known as Ada Commentaries and commonly referenced in the format AI-ddddd. For this implementation, the following tests were determined to be inapplicable for the reasons indicated; references to Ada Commentaries are included as appropriate.

IMPLEMENTATION DEPENDENCIES

The following 201 tests have floating-point type declarations requiring more digits than SYSTEM.MAX DIGITS:

```
C24113L..Y (14 tests) C35705L..Y (14 tests)
C35706L..Y (14 tests) C35707L..Y (14 tests)
C35708L..Y (14 tests) C35802L..Z (15 tests)
C45241L..Y (14 tests) C45321L..Y (14 tests)
C45421L..Y (14 tests) C45521L..Z (15 tests)
C45524L..Z (15 tests) C45621L..Z (15 tests)
C45641L..Y (14 tests) C46012L..Z (15 tests)
```

The following 21 tests check for the predefined type LONG INTEGER:

C354G4C	C45231C	C45304C	C45411C	C45412C
C45502C	C45503C	C45504C	C45504F	C45611C
C45612C	C45613C	C45614C	C45631C	C45632C
B52004D	C55B07A	B55B09C	B86001W	C86006C
CD7101F				

 ${\tt C35702B},\ {\tt C35713C},\ {\tt B86001U},\ {\tt and}\ {\tt C86006G}\ {\tt check}\ {\tt for}\ {\tt the}\ {\tt predefined}\ {\tt type}\ {\tt LONG}\ {\tt FLOAT}.$

C35713D and B86001Z check for a predefined floating-point type with a name other than FLOAT, LONG FLOAT, or SHORT FLOAT.

A35801E checks that FLOAT'FIRST..FLOAT'LAST may be used as a range constraint in a floating-point type declaration; for this implementation that range exceeds the safe numbers and must be rejected. (See section 2.3)

C45531M..P (4 tests) and C45532M..P (4 tests) check fixed-point operations for types that require a SYSTEM.MAX_MANTISSA of 47 or greater.

C45624A..B (2 tests) check that the proper exception is raised if MACHINE_OVERFLOWS is FALSE for floating point types; for this implementation, MACHINE OVERFLOWS is TRUE.

C86001F recompiles package SYSTEM, making package TEXT_IO, and hence package REPORT, obsolete. For this implementation, the package TEXT_IO is dependent upon package SYSTEM.

B86001Y checks for a predefined fixed-point type other than DURATION.

C96005B check: for values of type DURATION'BASE that are outside the range of DURATION. There are no such values for this implementation.

CD1009C uses a representation clause specifying a non-default size for a floating-point type.

CD2A84A, CD2A84E, CD2A84I...J (2 tests), and CD2A84O use representation clauses specifying non-default sizes for access types.

IMPLEMENTATION DEPENDENCIES

The tests listed in the following table are not applicable because the given file operations are not supported for the given combination of mode and file access method:

Test	File Operati	on Mode	File Access Method
CE2102D	CREATE	IN FILE	SEQUENTIAL IO
CE2102E	CREATE	OUT FILE	SEQUENTIAL 10
CE2102F	CREATE	INOUT FILE	DIRECT IO
CE2102I	CREATE	IN FILE	DIRECT ⁻ IO
CE2102J	CREATE	OUT FILE	DIRECTIO
CE2102N	OPEN	IN FILE	SEQUENTIAL IO
CE21020	RESET	IN FILE	SEQUENTIAL 10
CE2102P	OPEN	$\mathtt{OUT}_\mathtt{FILE}$	SEQUENTIAL_IO
CE2102Q	RESET	OUT_FILE	SEQUENTIAL_IO
CE2102R	OPEN	INOUT_FILE	
CE2102S	RESET	INOUT_FILE	DIRECT_10
CE2102T	OPEN	IN_FILE	DIRECT_IO
CE2102U	RESET	IN_FILE	DIRECT_IO
CE2102V	OPEN	OUT_FILE	DIRECT_IO
CE2102W	RESET	OUT_FILE	DIRECT_IO
CE3102E	CREATE	IN_FILE	TEXT_IO
CE3102F	RESET	An y Mode	TEXT_IO
CE3102G	DELETE		TEXT_IO
CE3102I	CREATE	OUT_FILE	TEXT_IO
CE3102J	OPEN	IN_FILE	TEXT_10
CE3102K	OPEN	OUT_FILE	TEXT_IO

CE2107A..E (5 tests), CE2107L, CE2110B, and CE2111D attempt to associate multiple internal files with the same external file for sequential files. The proper exception is raised when multiple access is attempted.

CE2107F..H (3 tests), CE2110D, and CE2111H attempt to associate multiple internal files with the same external file for direct files. The proper exception is raised when multiple access is attempted.

CE3111A..B (2 tests), CE3111D..E (2 tests), CE3114B, and CE3115A attempt to associate multiple internal files with the same external file for text files. The proper exception is raised when multiple access is attempted.

CE2203A checks that WRITE raises USE ERROR if the capacity of the external file is exceeded for SEQUENTIAL IO. This implementation does not restrict file capacity.

CE2403A checks that WRITE raises USE_ERROR if the capacity of the external file is exceeded for DIRECT_IO. This implementation does not restrict file capacity.

IMPLEMENTATION DEPENDENCIES

CE3304A checks that USE ERROR is raised if a call to SET LINE LENGTH or SET PAGE LENGTH specifies a value that is inappropriate for the external file. This implementation does not have inappropriate values for either line length or page length.

CE3413B checks that PAGE raises LAYOUT ERROR when the value of the page number exceeds COUNT'LAST. For this implementation, the value of COUNT'LAST is greater than 150000 making the checking of this objective impractical.

2.3 TEST MODIFICATIONS

Modifications (see section 1.3) were required for 22 tests.

The following tests were split into two or more tests because this implementation did not report the violations of the Ada Standard in the way expected by the original tests:

B24009A	B33301B	B38003A	B38003B	B38009A
B38009B	B85008G	B85008H	BC1303F	BC3005B
BD2BO3A	BD2D03A	BD4003A		

A35801E was graded inapplicable by Evaluation Modification as directed by the AVO; the compiler rejects the use of the range FLOAT'FIRST..FLOAT'LAST as the range constraint of a floating-point type declaration because the bounds lie outside of the range of safe numbers (cf. ARM 3.5.7(12)).

CD1009A, CD1009I, CD1C03A, CD2A22J, CD2A24A, and CD2A31A...C (3 tests) were graded passed by Evaluation Modification as directed by the AVO. These tests use instantiations of the support procedure LENGTH_CHECK, which uses Unchecked Conversion according to the interpretation given in AI-00590. The AVO ruled that this interpretation is not binding under ACVC 1.11; the tests are ruled to be passed if they produce Failed messages only from the instances of LENGTH_CHECK--i.e, the allowed Report.Failed messages have the general form:

[&]quot; * CHECK ON REPRESENTATION FOR <TYPE ID> FAILED."

APPENDIX C

APPENDIX F OF THE Ada STANDARD

The only allowed implementation dependencies correspond to implementation-dependent pragmas, to certain machine-dependent conventions as mentioned in Chapter 13 of the Ada Standard, and to certain allowed restrictions on representation clauses. The implementation-dependent characteristics of this Ada implementation, as described in this Appendix, are provided by the customer. Unless specifically noted otherwise, references in this Appendix are to compiler documentation and not to this report. Implementation-specific portions of the package STANDARD, which are not a part of Appendix F, are:

CHAPTER 3

PROCESSING INFORMATION

3.1 TESTING ENVIRONMENT

The host and target computer systems for the Ada implementation were connected by Ethernet. The rest of this Ada implementation tested in this validation effort is described adequately by the information given in the initial pages of this report.

For a point of contact for technical information about this Ada implementation system, see:

Steve Hodges Verdix Corporation 14130-A Sully Field Circle Chantilly VA 22021

For a point of contact for sales information about this Ada implementation system, see:

Steve Hodges Verdix Corporation 14130-A Sully Field Circle Chantilly VA 22021

Testing of this Ada implementation was conducted at the customer's site by a validation team from the AVF.

3.2 SUMMARY OF TEST RESULTS

An Ada Implementation passes a given ACVC version if it processes each test of the customized test suite in accordance with the Ada Programming Language Standard, whether the test is applicable or inapplicable; otherwise, the Ada Implementation fails the ACVC [Pro90].

PROCESSING INFORMATION

For all processed tests (inapplicable and applicable), a result was obtained that conforms to the Ada Programming Language Standard.

a) Total Number of Applicable Tests	3797
b) Total Number of Withdrawn Tests	81
c) Processed Inapplicable Tests	91
d) Non-Processed I/O Tests	0
e) Non-Processed Floating-Point	
Precision Tests	201

- f) Total Number of Inapplicable Tests 292 (c+d+e)
- g) Total Number of Tests for ACVC 1.11 4170 (a+b+f)

All I/O tests of the test suite were processed because this implementation supports a file system. The above number of floating-point tests were not processed because they used floating-point precision exceeding that supported by the implementation. When this compiler was tested, the tests listed in section 2.1 had been withdrawn because of test errors.

3.3 TEST EXECUTION

Version 1.11 of the ACVC comprises 4170 tests. When this compiler was tested, the tests listed in section 2.1 had been withdrawn because of test errors. The AVF determined that 292 tests were inapplicable to this implementation. All inapplicable tests were processed during validation testing except for 201 executable tests that use floating-point precision exceeding that supported by the implementation. In addition, the modified tests mentioned in section 2.3 were also processed.

A magnetic tape containing the customized test suite (see section 1.3) was taken on-site by the validation team for processing. The contents of the magnetic tape were not loaded directly onto the host computer. The tape was loaded onto a Sun Workstation, and the tests were copied over Ethernet to the host machine.

After the test files were loaded onto the host computer, the full set of tests was processed by the Ada implementation.

The tests were compiled and linked on the host computer system, as appropriate. The executable images were transferred to the target computer system by the communications link described above, and run. The results were captured on the host computer system.

Testing was performed using command scripts provided by the customer and reviewed by the validation team. See Appendix B for a complete listing of the processing options for this implementation. It also indicates the default options.

PROCESSING INFORMATION

Test output, compiler and linker listings, and job logs were captured on magnetic tape and archived at the AVF. The listings examined on-site by the validation team were also archived.

APPENDIX A

MACRO PARAMETERS

This appendix contains the macro parameters used for customizing the ACVC. The meaning and purpose of these parameters are explained in [UG89]. The following macro parameters are defined in terms of the value V of \$MAX_IN_LEN which is the maximum input line length permitted for the tested implementation. For these parameters, Ada string expressions are given rather than the macro values themselves.

Macro Parameter	Macro Value
\$BIG_ID1	(1V-1 => 'A', V => '1')
\$BIG_ID2	(1V-1 => 'A', V => '2')
\$BIG_ID3	(1V/2 => 'A') & '3' & (1V-1-V/2 => 'A')
\$BIG_ID4	(1V/2 => 'A') & '4' & (1V-1-V/2 => 'A')
\$BIG_INT_LIT	(1V-3 => '0') & "298"
\$BIG_REAL_LIT	(1V-5 => '0') & "690.0"
\$BIG_STRING1	""' & (1V/2 => 'A') & '"'
\$BIG_STRING2	""' & (1V-1-V/2 => 'A') & '1' & '"'
\$BLANKS	(1V-20 => ' ')
\$MAX_LEN_INT_BASED_LI	TERAL "2:" & (1V-5 => '0') & "11:"
\$MAX_LEN_REAL_BASED_L	ITERAL "16:" & (1V-7 => '0') & "F.E:"
\$MAX_STRING_LITERAL	""' & (1V-2 => 'A') & '"'

MACRO PARAMETERS

The following table lists all of the other macro parameters and their respective values:

Macro Parameter	Macro Value
\$MAX_IN_LEN	499
\$ACC_SIZE	32
\$ALIGNMENT	4
\$COUNT_LAST	2_147_483_647
\$DEFAULT_MEM_SIZE	16_777_216
\$DEFAULT_STOR_UNIT	8
\$DEFAULT_SYS_NAME	I386
\$DELTA_DOC	0.0000000004656612873077392578125
\$ENTRY_ADDRESS	SYSTEM."+"(16#40#)
\$ENTRY_ADDRESS1	SYSTEM."+"(16#80#)
\$ENTRY_ADDRESS2	SYSTEM."+"(16#100#)
\$FIELD_LAST	2_147_483_647
\$FILE_TERMINATOR	, ,
\$FIXED_NAME	NO_SUCH_TYPE
\$FLOAT_NAME	NO_SUCH_TYPE
\$FORM_STRING	11 11
\$FORM_STRING2	"CANNOT_RESTRICT_FILE_CAPACITY"
\$GREATER_THAN_DURATIO	0N 100_000.0
\$GREATER_THAN_DURATIO	N BASE LAST 10_000_000.0
\$GREATER_THAN_FLOAT_E	BASE_LAST 1.8E+308
\$GREATER_THAN_FLOAT_S	SAFE_LARGE 5.0E+307

MACRO PARAMETERS

SGREATER THAN SHORT FLOAT SAFE LARGE 9.0E+37 SHIGH PRIORITY 99 SILLEGAL EXTERNAL FILE NAME1 "/illegal/file name/2}]%2102c.dat" SILLEGAL EXTERNAL FILE NAME2 "/illegal/file_name/CE2102C*.DAT" \$INAPPROPRIATE_LINE_LENGTH SINAPPROPRIATE PAGE LENGTH \$INCLUDE PRAGMA1 PRAGMA INCLUDE ("A28006D1.TST") \$INCLUDE PRAGMA2 PRAGMA INCLUDE ("B28006D1.TST") -2 147 483 648 \$INTEGER FIRST 2 147 483 647 SINTEGER LAST SINTEGER LAST PLUS 1 2 147 483 648 SINTERFACE LANGUAGE \$LESS THAN DURATION -100 000.0 SLESS THAN DURATION BASE FIRST $-1\overline{0}$ 000 000.0 \$LINE TERMINATOR ASCII.LF & ASCII.FF \$LOW PRIORITY **\$MACHINE CODE STATEMENT** CODE $O'(OP \Rightarrow NOP)$; \$MACHINE CODE TYPE CODE 0 \$MANTISSA DOC 31 15 \$MAX DIGITS \$MAX INT 2 147 483 647 \$MAX INT PLUS 1 2 147 483 648

SMIN INT

-2 147 483 648

MACRO PARAMETERS

\$NAME TINY_INTEGER

\$NAME_LIST I386

\$NAME SPECIFICATION1 VMS386\$DKB0:[ACVC1.11.C.E]X2120A.

\$NAME_SPECIFICATION2 VMS386\$DKB0:[ACVC1.11.C.E]X2120B.

\$NAME SPECIFICATION3 VMS386\$DKB0:[ACVCl.11.C.E]X3119A.

\$NEG_BASED_INT 16#F000000E#

\$NEW_MEM_SIZE 16_777_216

SNEW_STOR_UNIT 8

\$NEW_SYS_NAME I386

\$PAGE TERMINATOR ASCII.FF

\$RECORD DEFINITION RECORD SUBP: OPERAND; END RECORD;

\$RECORD NAME CODE 0

\$TASK TZE 32

\$TASK_STORAGE_SIZE 1024

\$TICK 0.01

\$VARIABLE_ADDRESS VAR 3'ADDRESS

\$VARIABLE ADDRESS1 VAR 1'ADDRESS

\$VARIABLE ADDRESS2 VAR 2'ADDRESS

SYOUR PRAGMA PASSIVE

APPENDIX B

COMPILATION SYSTEM OPTIONS

The compiler options of this Ada implementation, as described in this Appendix, are provided by the customer. Unless specifically noted otherwise, references in this appendix are to compiler documentation and not to this report.

VADS ADA - Ada compiler

Syntax

VADS ADA source file [, ...]

Qualifiers

/ALPEND Append all output to a log file.

/DEBUG

/DEBUG=G Write out the gnrx.lib file in ASCII.

/DEFINE

/DEFINE=(identifier:type=value", ...)

Define identifier of a specified type and value. See VADS ADA PREPROCESSOR REFERENCE.

/DEPENDENCIES Analyze for dependencies only; no link will be performed this option is given (/MAIN and /OUTPUT options must not be used with this qualifier).

/ERRORS

/ERRORS[=(option [, ...])] Process compilation error messages using the ERR tool and direct the output to SYS\$OUTPUT; the parentheses can be omitted if only one qualifier is given (by default, only lines containing errors are listed).

Options:

LISTING List entire input file.

EDITOR[="editor"]

Insert error messages into the source file and call a text editor (EDT by default). If a value is given as a quoted string, that string is used to invoke the editor. This allows other editors to be used instead of the default.

OUTPUT[=file name]

Direct error processed output to the specified file name; if no file name is given, the source file name is used with a file extension .ERR.

BRIEF List only the affected lines [default]

Only one of the BRIEF, LISTING, OUTPUT, or EDITOR options can be used in a single command.

For more information about the /ERRORS option, see also COMPILING ADA PROGRAMS, COMPILER ERROR MESSAGE PROCESSING.

/EXECUTABLE

/EXECUTABLE=file name

Provide an explicit name for the executable when used with the /MAIN qualif the file name value must be supplied (if the file type is omitted,

.EXE is assumed).

/KEEP_II. Keep the intermediate language (IL) file produced by the compiler front end. The IL file will be placed in the OBJECTS directory, with the name ADA_SOURCE.I.

/LIBRARY

/LIBRARY=library_name Operate in VADS library library name (the current working directory is the default).

/LINK ARGUMENTS

/LINK ARGUMENTS="value"

Pass command qualifiers and parameters to the linker.

/MAIN

/MAIN[=unit_name] Produce an executable program using the named unit as the main program; if no value is given, the name is derived from the first Ada file name parameter (the .A suffix is removed); the executable file name is derived from the main program name unless the /EXECUTABLE qualifier is used.

/NOOPTIMIZE

Do not optimize.

/WARNINGS Print warning diagnostics.

/OPTIMIZE

/OPTIMIZE[=number] Invoke the code optimizer (OPTIM3). An optional digit provides the level of optimization. /OPTIMIZE=4 is the default.

```
/OPTIMIZE no digit, full optimization
/OPTIMIZE=0
                  prevents optimization
/OPTIMIZE=1
                  no hoisting
/OPTIMIZE=2
                  no hoisting, but more passes
                  no hoisting, but even more passes
/OPTIMIZE=3
/OPTIMIZE≈4
                  hoisting from loops
/OPTIMIZE=5
                  hoisting from loops, but more passes
/OPTIMIZE=6
                  hoisting from loops with maximum passes
/OPTIMIZE=7
                  hoisting from loops and branches
/OPTIMIZE=8
                  hoisting from loops and branches, more passes
                  hoisting from loops and branches, maximum passes
/OPTIMIZE=9
```

Hoisting from branches (and cases alternatives) can be slow and does not always provide significant performance gains, so it can be suppressed.

For more information about optimization, see COMPILING ADA PROGRAMS, Optimization. See also pragma OPTIMIZE CODE(OFF).

/OUTPUT

/OUTPUT=file_name Direct the output to file_name (the default is SYS\$OUTPUT).

/PRE PROCESS Invoke the Ada Preprocessor. See VADS ADA PREPROCESSOR REFERENCE.

COMPILATION SYSTEM OPTIONS

/RECOMPILE_LIBRARY=VADS_library
Force analysis of all generic instantiations causing reinstantiation of any that are out of date.

/SHOW Show the name of the tool executable but do not execute it.

/SHOW ALL Print the name of the front end, code generator, optimizer, and linker, and list the tools that will be invoked.

/SUPPRESS Apply pragma SUPPRESS for all checks to the entire compilation. (See also pragma SUPPRESS(ALL CHECKS))

/TIMING Print timing information for the compilation.

/VERBOSE Print information for the compilation.

Description

The command VADS ADA executes the Ada compiler and compiles the named Ada source file, ending with the .A suffix. The file must reside in a VADS library directory. The ADA.LIB file in this directory is modified after each Ada unit is compiled.

By default, VADS ADA produces only object and net files. If the /MAIN option is used, the compiler automatically invokes VADS LD and builds a complete program with the named library unit as the main program.

Non-Ada object files may be given as arguments to VADS ADA. These files will be passed on to the linker and will be linked with the specified Ada object files.

Command line options may be specified in any order, but the order of compilation and the order of the files to be passed to the linker can be significant.

Several VADS compilers may be simultaneously available on a single system. The VADS ADA command within any version of VADS on a system will execute the correct compiler components based upon visible library directives.

Program listings with a disassembly of machine code instructions are generated by VADS DB or VADS DAS.

Diagnostics

The diagnostics produced by the VADS compiler are intended to be self-explanatory. Most refer to the RM. Each RM reference includes a section number and optionally, a paragraph number enclosed in parentheses.

L'NKER OPTIONS

The linker options of this Ada implementation, as described in this Appendix, are provided by the customer. Unless specifically noted otherwise, references in this appendix are to linker documentation and not to this report.

VADS LD - prelinker

Syntax

VADS LD unit name

Qualifiers

/AP IND

/DEBUG Debug memory overflow (use in cases where linking a large number units causes the error message "local symbol overflow" to occur).

/EARLY

/EARLY="unit_name" Force the given unit to be elaborated as early as possible (unit_name must be enclosed in double quotes).

/EXECUTABLE

/EXECUTABLE[=file name]

Put the output in the named file. The default executable names are <main_unit>.EXE on self-hosts or <main_unit>.VOX on cross targets.

/FILES Print a list of dependent files in elaboration order and suppres linking.

/LIBRARY

/LIBRARY=library_name Operate in VADS library_library_name (the current working directory is the default).

/LINK OPTIONS

/LINK_OPTIONS=object_file_or_qualifier [,...]"

Add the options surrounded by quotes to the invocation of the linker.

COMPILATION SYSTEM OPTIONS

/OUTPUT

/OUTPUT=file name Direct output to file name. Default is SYS\$OUTPUT.

/SHOW Show the name of the tool executable but do not execute it.

/UNITS Print a list of dependent units in order and suppress linking.

/VERBOSE Print the VMS linker command prior to execution.

/VERIFY Print the VMS linker command but suppress execution.

Description

VADS LD collects the object files needed to make unit name a main program and calls the VMS linker to link together all Ada and other language objects required to produce an executable. unit_name must be a non-generic subprogram that is either a procedure or a function that returns an Ada STANDARD.INTEGER (the predefined type INTEGER). The utility uses the net files produced by the Ada compiler to check dependency information. VADS LD produces an exception mapping table, a unit elaboration table, and passes this information to the linker.

VADS LD reads instructions for generating executables from the ADA.LIB file in the VADS libraries on the search list. Besides information generated by the compiler, these directives also include WITHn directives that allow the automatic linking of object modules compiled from other languages or Ada object modules not named in context clauses in the Ada source. Any number of WITH directives may be placed into a library, but they must be numbered contiguously beginning at WITH1. The directives are recorded in the library's ADA.LIB file and have the following form:

WITH1|LINK|object_file| WITH2|LINK|archive_file|

WITH directives may be placed in the local Ada libraries or in any VADS library on the search list.

A WITH directive in a local VADS library or earlier on the library search list will hide the same numbered WITH directive in a library later in the library search list.

Use VADS INFO to change or report library directives in the current library.

All arguments after unit_name are passed on to the linker. These arguments may be linker options, names of object files or archive libraries, or library abbreviations.

Diagnostics

Self-explanatory diagnostics are produced for missing files, etc. Occasional additional messages are produced by the linker.

COMPILATION SYSTEM OPTIONS

Files

Normally VADS LD generates an intermediate file with the process ID as a substring, VADSOPTIONcess_ID>.OPT.

With either the /VERIFY or /VERBOSE qualifiers,; however, VADS LD will produce the intermediate file <main_unit>.OPT.

ATTACHMENT I

APPENDIX F. Implementation-Dependent Characteristics

1. Implementation-Dependent Pragmas

1.1. INITIALIZE Pragma

Takes one of the identifiers STATIC or DYNAMIC as the single argument. This pragma is only allowed within a library-level package spec or body. It specifies that all objects in the package be initialized as requested by the pragma (i.e statically or dynamically). Only library-level objects are subject to static initialization; all objects within procedures are always (by definition) dynamic. If pragma INITIALIZE(STATIC) is used, and an object cannot be initialized statically, code will be generated to initialize the object, and a warning message will be generated.

1.2. INLINE ONLY Pragma

The INLINE ONLY pragma, when used in the same way as pragma INLINE, indicates to the compiler that the subprogram must always be inlined. This pragma also suppresses the generation of a callable version of the routine which saves code space. If a user erroneously makes an INLINE ONLY subprogram recursive, a warning message will be emitted and an PROGRAM ERROR will be raised at run time.

1.3. BUILT_IN Pragma

The BUILT IN pragma is used in the implementation of some predefined Ada packages, but provides no user access. It is used only to implement code bodies for which no actual Ada body can be provided, for example, the MACHINE_CODE package.

1.4. SHARE_CODE Pragma

The SHARE CODE pragma takes the name of a generic instantiation or a generic unit as the first argument and one of the identifiers TRUE or FALSE as the second argument. This pragma is only allowed immediately at the place of a

declarative item in a declarative part or package specification, or after a library unit in a compilation, but before any subsequent compilation unit.

When the first argument is a generic unit, the pragma applies to all instantiations of that generic. When the first argument is the name of a generic instantiation, the pragma applies only to the specified instantiation, or overloaded instantiations.

If the second argument is TRUE, the compiler will try to share code generated for a generic instantiation with code generated for other instantiations of the same generic. When the second argument is FALSE, each instantiation will get a unique copy of the generated code. The extent to which code is shared between instantiations depends on this pragma and the kind of generic formal parameters declared for the generic unit.

The name pragma SHARE_BODY is also recognized by the implementation and has the same effect as SHARE_CODE. It is included for compatability with earlier versions of VADS.

1.5. NO IMAGE Pragma

The pragma suppresses the generation of the image array used for the IMAGE attribute of enumeration types. This eliminates the overhead required to store the array in the executable image. An attempt to use the IMAGE attribute on a type whose image array has been suppressed will result in a compilation warning and PROGRAM_ERROR raised at run time.

1.6. EXTERNAL NAME Pragma

The EXTERNAL NAME pragma takes the name of a subprogram or variable defined in Ada and allows the user to specify a different external name that may be used to reference the entity from other languages. The pragma is allowed at the place of a declarative item in a package specification and must apply to an object declared earlier in the same package specification.

1.7. INTERFACE NAME Pragma

The INTERFACE NAME pragma takes the name of a a variable or subprogram defined in another language and allows it to be referenced directly in Ada. The pragma will replace all occurrences of the variable or subprogram name with an external reference to the second, link argument. The pragma is allowed at the place of a declarative item in a package specification and must apply to an object or subprogram declared earlier in the same package specification. The

APPENDIX F OF THE Ada ST'NDARD

object must be declared as a scalar or an access type. The object cannot be any of the following:

a loop variable, a constant, an initialized variable, an array, or a record.

1.8. IMPLICIT CODE Pragma

Takes one of the identifiers ON or OFF as the single argument. This pragma is only allowed within a machine code procedure. It specifies that implicit code generated by the compiler be allowed or disallowed. A warning is issued if OFF is used and any implicit code needs to be generated. The default is ON.

1.9. OPTIMIZE CODE Pragma

Takes one of the identifiers ON or OFF as the single argument. This pragma is only allowed within a machine code procedure. It specifies whether the code should be optimized by the compiler. The default is ON. When OFF is specified, the compiler will generate the code as specified.

2. Implementation of Predefined Pragmas

2.1. CONTROLLED

This pragma is recognized by the implementation but has no effect.

2.2. ELABORATE

This pragma is implemented as described in Appendix B of the Ada RM.

2.3. INLINE

This pragma is implemented as described in Appendix B of the Ada RM.

2.4. INTERFACE

This pragma supports calls to 'C' and FORTRAN functions. The Ada subprograms can be either functions or procedures. The types of parameters and the result type for functions must be scalar, access, or the predefined type ADDRESS in SYSTEM. All parameters must have mode IN. Record and array objects can be passed by reference using the ADDRESS attribute.

2.5. LIST

This pragma is implemented as described in Appendix B of the Ada RM.

2.6. MEMORY_SIZE

This pragma is recognized by the implementation. The implementation does not allow SYSTEM to be modified by means of pragmas; the SYSTEM package must be recompiled.

2.7. NON REENTRANT

This pragma takes one argument which can be the name of either a library subprogram or a subprogram declared immediately within a library package spec or body. It indicates to the compiler that the subprogram will not be called recursively, allowing the compiler to perform specific optimizations. The pragma can be applied to a subprogram or a set of overloaded subprograms within a package spec or package body.

2.8. NOT ELABORATED

This pragma can only appear in a library package specification. It indicates that the package will not be elaborated because it is either part of the RTS, a configuration package, or an Ada package that is referenced from a language other than Ada. The presence of this pragma suppresses the generation of elaboration code and issues warnings if elaboration code is required.

2.9. OPTIMIZE

This pragma is recognized by the implementation but has no effect.

2.10. PACK

This pragma will cause the compiler to choose a non-aligned representation for composite types. It will not cause objects to be packed at the bit level.

2.11. PAGE

This pragma is implemented as described in Appendix B of the Ada RM.

2.12. PASSIVE

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The pragma has three forms:

PRAGMA PASSIVE; PRAGMA PASSIVE(SEMAPHORE); PRAGMA PASSIVE(INTERRUPT, <number>);

This pragma Pragma passive can be applied to a task or task type declared immediately within a library package spec or body. The pragma directs the compiler to optimize certain tasking operations. It is possible that the statements in a task body will prevent the intended optimization; in these cases, a warning will be generated at compile time and will raise TASKING ERROR at runtime.

2.13. PRIORITY

This pragma is implemented as described in Appendix B of the Ada RM.

2.14. SHARED

This pragma is recognized by the implementation but has no effect.

2.15. STORAGE UNIT

This pragma is recognized by the implementation. The implementation does not allow SYSTEM to be modified by means of pragmas; the SYSTEM package must be recompiled.

2.16. SUPPRESS

This pragma is implemented as described, except that DIVISION_CHECK and in some cases OVERFLOW_CHECK cannot be suppressed.

2.17. SYSTEM NAME

This pragma is recognized by the implementation. The implementation does not allow SYSTEM to be modified by means of pragmas; the SYSTEM package must be recompiled.

3. Implementation-Dependent Attributes

3.1. P'REF

For a prefix that denotes an object, a program unit, a label, or an entry:

This attribute denotes the effective address of the first of the storage units allocated to P. For a subprogram, package, task unit, or label, it refers to the address of the machine code associated with the corresponding body or statement. For an entry for which an address clause has been given, it refers to the corresponding hardware interrupt. The attribute is of the type OPERAND defined in the package MACHINE_CODE. The attribute is only allowed within a machine code procedure.

See section F.4.8 for more information on the use of this attribute.

(For a package, task unit, or entry, the 'REF attribute is not supported.)

3.2. T'TASKID

For a task object or a value T, T'TASK_ID yields the unique task id associated with a task. The value of this attribute is of the type ADDRESS in the package SYSTEM.

- 4. Specification of Package SYSTEM
- -- Copyright 1987, 1988, 1989, 1990 Verdix Corporation

with unsigned types; package SYSTEM is

```
pragma suppress(ALL_CHECKS);
pragma suppress(EXCEPTION_TABLES);
pragma not elaborated;
```

type NAME is (i386);

SYSTEM NAME := i386;

STORAGE UNIT : constant := 8;

MEMORY SIZE : constant := 16 777 216;

-- System-Dependent Named Numbers

MIN INT : constant := -2 147 483 648; MAX INT : constant := 2 147 483 647;

MAX DIGITS : constant := 15;

MAX MANTISSA : constant := 31;

FINE DELTA : constant := 2.0**(-31);

TICK : constant := 0.01;

-- Other System-dependent Declarations

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```
subtype PRIORITY is INTEGER range 0 .. 99;
        MAX REC SIZE : integer := 64*1024;
        type ADDRESS is private;
        function ">" (A: ADDRESS; B: ADDRESS) return BOOLEAN;
        function "<" (A: ADDRESS; B: ADDRESS) return BOOLEAN;
        function ">="(A: ADDRESS; B: ADDRESS) return BOOLEAN;
        function "<="(A: ADDRESS; B: ADDRESS) return BOOLEAN;
        function "-" (A: ADDRESS; B: ADDRESS) return INTEGER;
        function "+" (A: ADDRESS; I: INTEGER) return ADDRESS;
        function "-" (A: ADDRESS; I: INTEGER) return ADDRESS;
        function "+" (I: UNSIGNED TYPES.UNSIGNED INTEGER) return ADDRESS;
        function MEMORY ADDRESS
                 (I: UNSIGNED TYPES.UNSIGNED INTEGER) return ADDRESS renames "+";
        NO ADDR : constant ADDRESS;
        type TASK ID is private;
        NO TASK ID: constant TASK ID;
        type PROGRAM ID is private;
        NO PROGRAM I\overline{D}: constant PROGRAM ID;
private
        type ADDRESS is new UNSIGNED TYPES.UNSIGNED INTEGER;
        NO ADDR : constant ADDRESS := 0;
        pragma BUILT_IN(">");
        pragma BUILT_IN("<");</pre>
        pragma BUILT IN(">=");
        pragma BUILT IN("<=");</pre>
        pragma BUILT IN("-");
        pragma BUILT IN("+");
        type TASK ID is new UNSIGNED TYPES.UNSIGNED INTEGER;
        NO TASK I\overline{D}: constant TASK I\overline{D}:= 0;
        type PROGRAM ID is new UNSIGNED TYPES.UNSIGNED INTEGER;
        NO PROGRAM \overline{ID}: constant PROGRAM \overline{ID}:= 0;
end SYSTEM;
```

5. Restrictions on Representation Clauses

5.1. Pragma PACK

In the absence of Pragma PACK, record components are padded so as to provide for efficient access by the target hardware; pragma PACK applied to a record eliminates the padding where possible. Pragma PACK has no other effect on the storage allocated for record components a record representation is required.

5.2. Size Clauses

For scalar types, a representation clause will pack to the number of bits required to represent the range of the subtype. A size clause applied to a record type will not cause packing of components; an explicit record representation clause must be given to specify the packing of the components. A size clause applied to a record type will cause packing of components only when the component type is a discrete type. An error will be issued if there is insufficient space allocated. The SIZE attribute is not supported for task, access, or floating point types.

5.3. Address Clauses

Address clauses are only supported for variables. Since default initialization of a variable requires evaluation of the variable address, elaboration ordering requirements prohibit initialization of variables which have address clauses. The specified address indicates the physical address associated with the variable.

5.4. Interrupts

Interrupt entries are not supported.

5.5. Representation Attributes

The ADDRESS attribute is not supported for the following entities:

Packages Tasks Labels Entries

5.6. Machine Code Insertions

Machine code insertions are supported.

The general definition of the package MACHINE CODE provides

an assembly language interface for the target machine. It provides the necessary record type(s) needed in the code statement, an enumeration type of all the opcode mnemonics, a set of register definitions, and a set of addressing mode functions.

The general syntax of a machine code statement is as follows:

where n indicates the number of operands in the aggregate.

A special case arises for a variable number of operands. The operands are listed within a subaggregate. The format is as follows:

For those opcodes that require no operands, named notation must be used (cf. RM 4.3(4)).

The opcode must be an enumeration literal (i.e., it cannot be an object, attribute, or a rename).

An operand can only be an entity defined in MACHINE_CODE or the 'REF attribute.

The arguments to any of the functions defined in MACHINE CODE must be static expressions, string literals, or the functions defined in MACHINE CODE. The 'REF attribute may not be used as an argument in any of these functions.

Inline expansion of machine code procedures is supported.

6. Conventions for Implementation-generated Names

There are no implementation-generated names.

7. Interpretation of Expressions in Address Clauses

Address expressions in an address clause are interpreted as physical addresses.

8. Restrictions on Unchecked Conversions

None.

9. Restrictions on Unchecked Deallocations

None.

10. Implementation Characteristics of I/O Packages

Instantiations of DIRECT_IO use the value MAX_REC_SIZE as the record size (expressed in STORAGE_UNITS) when the size of ELEMENT_TYPE exceeds that value. For example, for unconstrained arrays such as string, where ELEMENT_TYPE'SIZE is very large, MAX_REC_SIZE is used instead. MAX_RECORD_SIZE is defined in SYSTEM and can be changed by a program before instantiating DIRECT_IO to provide an upper limit on the record size. In any case, the maximum size supported is 1024 x 1024 x STORAGE_UNIT_bits. DIRECT_IO will raise USE_ERROR if MAX_REC_SIZE_exceeds this absolute limit.

Instantiations of SEQUENTIAL IO use the value MAX REC SIZE as the record size (expressed in STORAGE_UNITS) when the size of ELEMENT_TYPE exceeds that value. For example, for unconstrained arrays such as string, where ELEMENT_TYPE'SIZE is very large, MAX_REC_SIZE is used instead. MAX_RECORD_SIZE is defined in SYSTEM and can be changed by a program before instantiating INTEGER_IO to provide an upper limit on the record size. SEQUENTIAL_IO imposes no limit on MAX_REC_SIZE.

11. Implementation Limits

The following limits are actually enforced by the implementation. It is not intended to imply that resources up to or even near these limits are available to every program.

11.1. Line Length

The implementation supports a maximum line length of 500 characters including the end of line character.

11.2. Record and Array Sizes

The maximum size of a statically sized array type is 4,000,000 x STORAGE_UNITS. The maximum size of a statically sized record type is 4,000,000 x STORAGE_UNITS. A record type or array type declaration that exceeds these limits will generate a warning message.

11.3. Default Stack Size for Tasks

In the absence of an explicit STORAGE_SIZE length specification, every task except the main program is allocated a fixed size stack of 10,240 STORAGE_UNITS. This is the value returned by T'STORAGE_SIZE for a task type T.

11.4. Default Collection Size

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In the absence of an explicit STORAGE_SIZE length attribute, the default collection size for an access type is 100 times the size of the designated type. This is the value returned by T'STORAGE_SIZE for an access type T.

11.5. Limit on Declared Objects

There is an absolute limit of 6,000,000 x STORAGE_UNITS for objects declared statically within a compilation unit. If this value is exceeded, the compiler will terminate the compilation of the unit with a FATAL error message.